1. Intro. This program generates DLX3 data that finds all "motley dissections" of an $m \times n$ rectangle into subrectangles.

The allowable subrectangles $[a\,..\,b) \times [c\,..\,d)$ have $0 \le a < b \le m$, $0 \le c < d \le n$, with $(a,b) \ne (0,m)$ and $(c,d) \ne (0,n)$; so there are $\bigl(\binom{m+1}{2}-1\bigr)\cdot\bigl(\binom{n+1}{2}-1\bigr)$ possibilities. Such a dissection is *motley* if the pairs (a,b) are distinct, and so are the pairs (c,d); in other words, no two subrectangles have identical top-bottom boundaries or left-right boundaries.

Furthermore we require that every $x \in [0\,..\,m)$ occurs at least once among the a's; every $y \in [0\,..\,n)$ occurs at least once among the c's. Otherwise the dissection could be collapsed into a smaller one, by leaving out that coordinate value.

It turns out that we can save a factor of (roughly) 2 by using symmetry, and looking at the unique rectangles that lie within the top and bottom rows of every solution.

```
#define maxd 36   /* maximum value for m or n */
#define encode(v) ((v) < 10 ? (v) + '0' : (v) - 10 + 'a')   /* encoding for values < 36 */
#include <stdio.h>
#include <stdlib.h>
  int m, n;   /* command-line parameters */
  main(int argc, char *argv[])
  {
    register int a, b, c, d, j, k;
    ⟨Process the command line 2⟩;
    ⟨Output the first line 3⟩;
    for (a = 0; a < m; a++)
      for (b = a + 1; b ≤ m; b++)
        if (a ≠ 0 ∨ b ≠ m) {
          for (c = 0; c < n; c++)
            for (d = c + 1; d ≤ n; d++)
              if (c ≠ 0 ∨ d ≠ n) {⟨Output the line for [a..b] × [c..d] 5⟩}
        }
  }
```

2. ⟨Process the command line 2⟩ ≡

```
  if (argc ≠ 3 ∨ sscanf(argv[1], "%d", &m) ≠ 1 ∨ sscanf(argv[2], "%d", &n) ≠ 1) {
    fprintf(stderr, "Usage: %s m n\n", argv[0]);
    exit(−1);
  }
  if (m > maxd ∨ n > maxd) {
    fprintf(stderr, "Sorry, m and n must be at most %d!\n", maxd);
    exit(−2);
  }
  printf("| motley-dlx %d %d\n", m, n);
```

This code is used in section 1.

3. The main primary columns jk ensure that cell (j,k) is covered, for $0 \le j < m$ and $0 \le k < n$. We also have secondary columns xab and ycd, to ensure that no interval is repeated. And there are primary columns xa and yc for the at-least-once conditions.

⟨ Output the first line 3 ⟩ ≡

```
for (j = 0; j < m; j++)
   for (k = 0; k < n; k++) printf("␣%c%c", encode(j), encode(k));
for (a = 1; a < m; a++) printf("␣1:%d|x%c", m − a, encode(a));
for (c = 1; c < n; c++) printf("␣1:%d|y%c", n − c, encode(c));
printf("␣|");
for (a = 0; a < m; a++)
   for (b = a + 1; b ≤ m; b++)
      if (a ≠ 0 ∨ b ≠ m) printf("␣x%c%c", encode(a), encode(b));
for (c = 0; c < n; c++)
   for (d = c + 1; d ≤ n; d++)
      if (c ≠ 0 ∨ d ≠ n) printf("␣y%c%c", encode(c), encode(d));
⟨ Output also the secondary columns for symmetry breaking 6 ⟩;
printf("\n");
```

This code is used in section 1.

4. Now let's look closely at the problem of breaking symmetry. For concreteness, let's suppose that $m = 7$ and $n = 8$. Every solution will have exactly one entry with interval x67, specifying a rectangle in the bottom row (since $m - 1 = 6$). If that rectangle has y57, say, a left-right reflection would produce an equivalent solution with y13; therefore we do not allow the rectangle for which $(a,b,c,d) = (6,7,5,7)$. Similarly we disallow $(6,7,c,d)$ whenever $8 - d < c$, since we'll find all solutions with $(6,7,8-d,8-c)$ that are left-right reflections of the solutions excluded.

If $a = 6$, $b = 7$, and $c + d = 8$, left-right reflection doesn't affect the rectangle in the bottom row. But we can still break the symmetry by looking at the top row, the rectangle whose specifications (a',b',c',d') have $(a',b') = (0,1)$. Let's introduce secondary columns !1, !2, !3, using !c when $c + d = 8$ at the bottom. Then if we put !1, !2, and !3 on every top-row rectangle with $c' + d' > 8$, we'll forbid such rectangles whenever the bottom-row policy has not already broken left-right symmetry. Furthermore, when $c' + d' = 8$ at the top, we put !1 together with x01 y26, and we put both !1 and !2 together with x01 y35. It can be seen that this completely breaks left-symmetry in all cases, because no solution has $c = c'$ and $d = d'$.

(Think about it.)

It's tempting to believe, as the author once did, that the same idea will break top-bottom symmetry too. But that would be fallacious: Once we've fixed attention on the bottommost row while breaking left-right symmetry, we no longer have any symmetry between top and bottom.

(Think about that, too.)

5. ⟨Output the line for $[a\,..\,b]\times[c\,..\,d]$ 5⟩ ≡

```
if (a ≡ m − 1 ∧ c + d > n) continue;    /* forbid this case */
for (j = a; j < b; j++)
  for (k = c; k < d; k++) printf("␣%c%c", encode(j), encode(k));
if (a ≡ m − 1 ∧ c + d ≡ n) printf("␣!%d", c);    /* flag a symmetric bottom row */
if (b ≡ 1 ∧ c + d ≥ n) {    /* disallow top rectangle if bottom one is symmetric */
  if (c + d ≠ n)
    for (k = 1; k + k < n; k++) printf("␣!%d", k);
  else
    for (k = 1; k < c; k++) printf("␣!%d", k);
}
if (a) printf("␣x%c", encode(a));
if (c) printf("␣y%c", encode(c));
printf("␣x%c%c␣y%c%c\n", encode(a), encode(b), encode(c), encode(d));
```

This code is used in section 1.

6. ⟨Output also the secondary columns for symmetry breaking 6⟩ ≡

```
for (k = 1; k + k < n; k++) printf("␣!%d", k);
```

This code is used in section 3.

7. Index.

MOTLEY-DLX